

# Curvature and Complexity:

Better lower bounds for geodesically convex optimization

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Chris Criscitiello

Nicolas Boumal

OPTIM, Chair of Continuous Optimization

Institute of Mathematics, EPFL



# Setting: $g$ -convex optimization

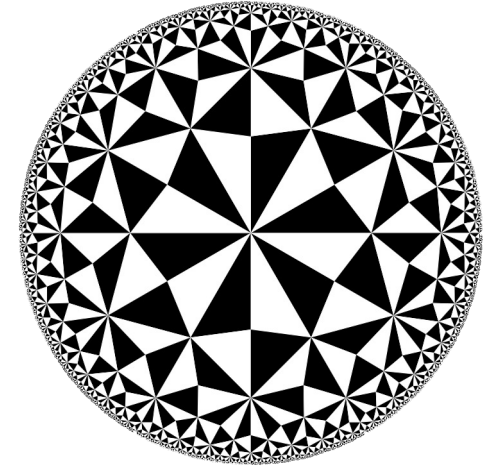
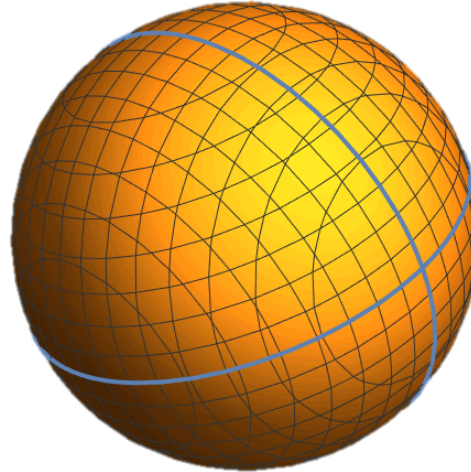
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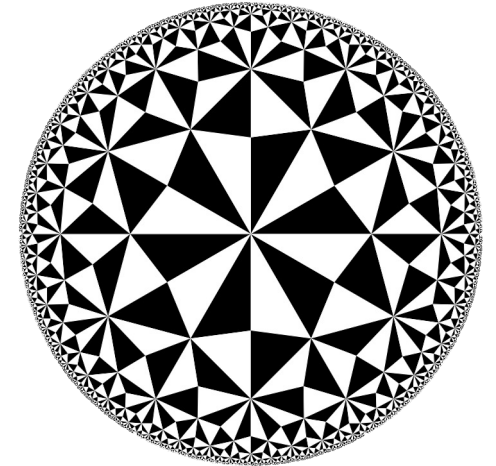
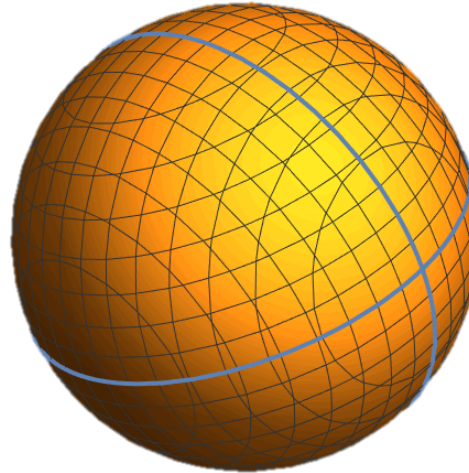
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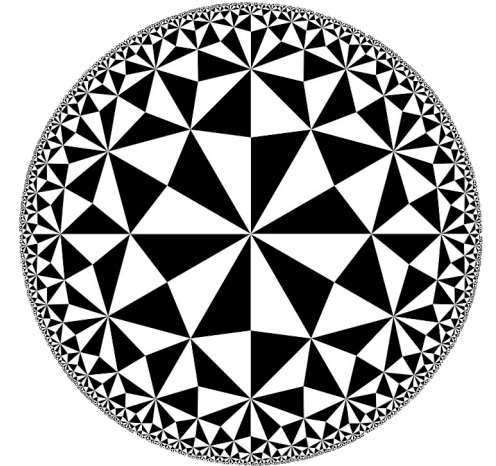
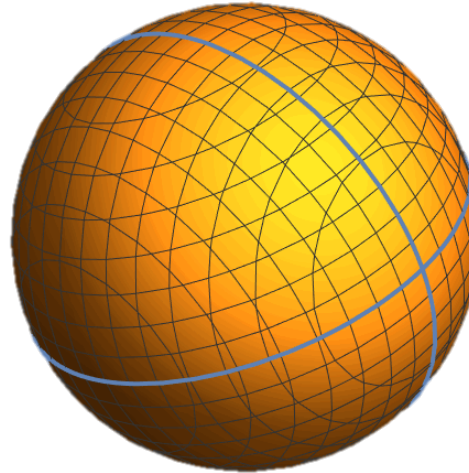
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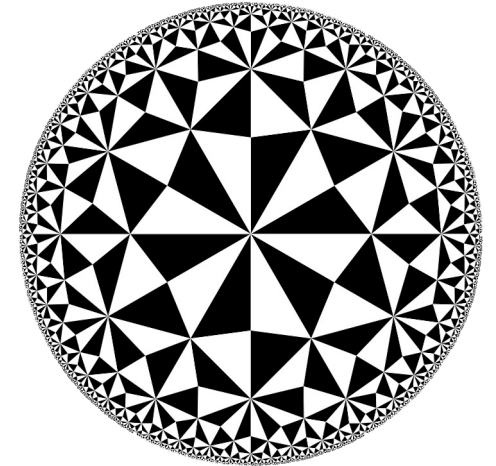
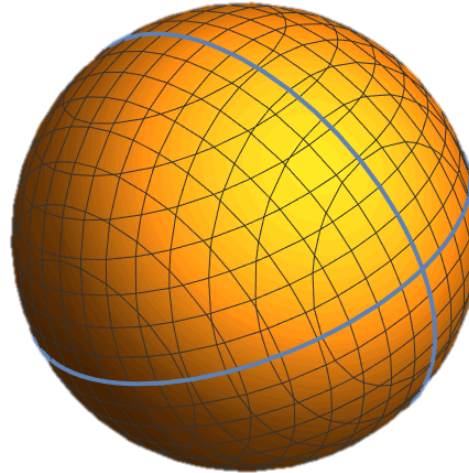
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$f$  is geodesic

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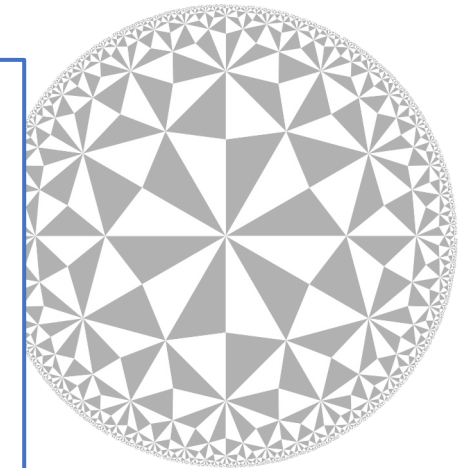
function

Most positively curved spaces do not carry nonconstant global g-convex functions.

Applications of g-convexity usually have  $K \leq 0$ :

- *Operator scaling*
- *Robust covariance estimation*
- *Matrix normal models*
- *Intrinsic medians (e.g., for phylogenetics)*

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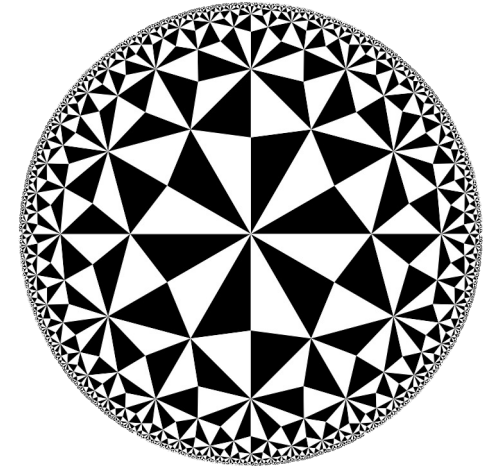
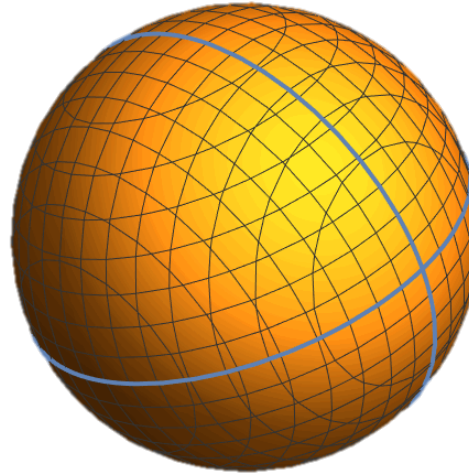
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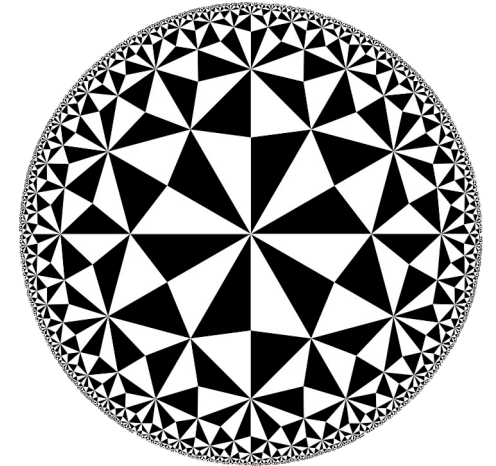
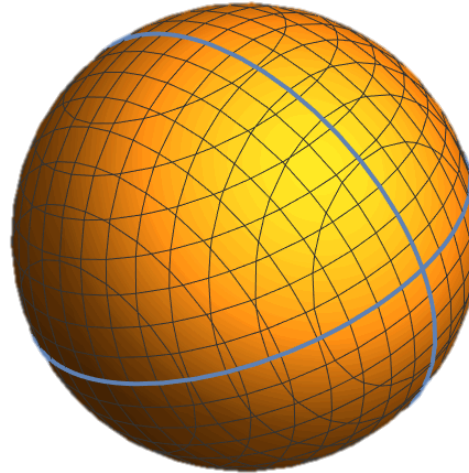
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Assume  $\mathcal{M}$  has constant curvature (simplifying assumption)

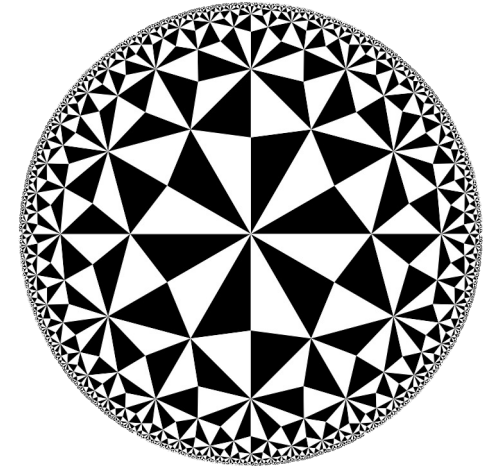
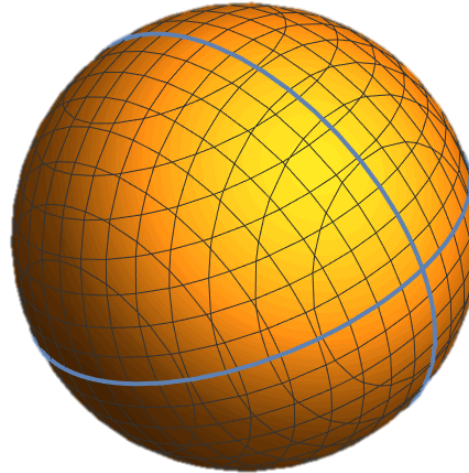
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Throughout  $\mathcal{M} = \mathbb{H}^d$  is a hyperbolic space of curvature  $K = -1$   
(simplifying assumption)

# Curvature

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$$K < 0$$

*Constant  
curvature:*

Euclidean space  $\mathbb{R}^d$

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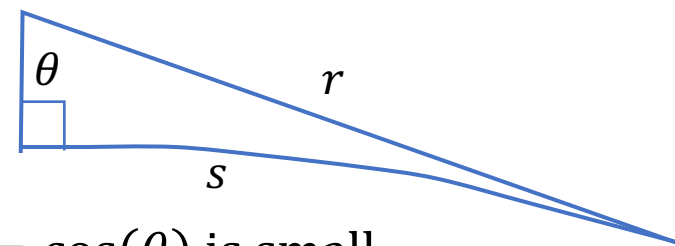
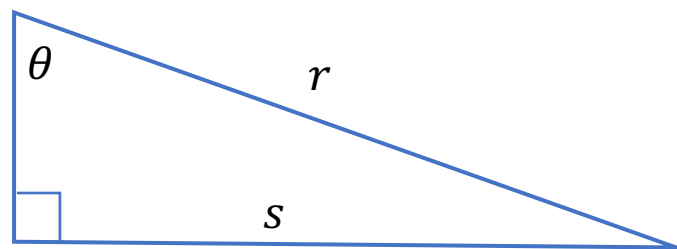
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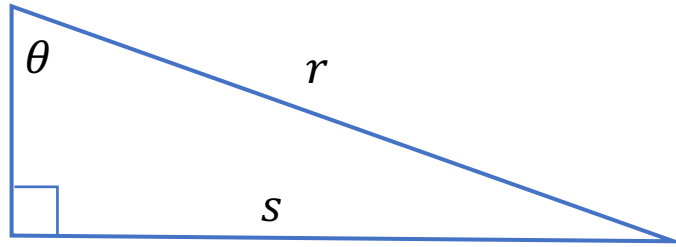
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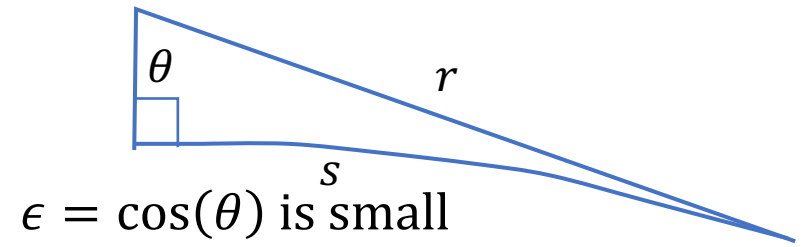


$\epsilon = \cos(\theta)$  is small

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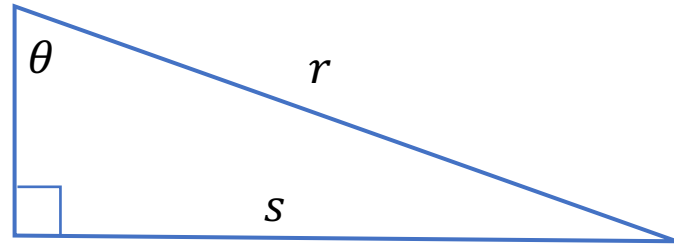
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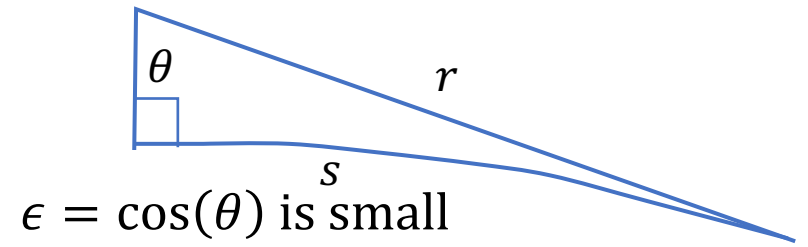
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$$\zeta = \frac{r\sqrt{|K|}}{\tanh(r\sqrt{|K|})} \sim 1 + r\sqrt{|K|}$$

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- (P3) **Smooth**  $|\nabla f(x) - P_{y \rightarrow x} \nabla f(y)| \leq L \text{dist}(x, y)$
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Worst-case complexity for (P2) & (P3): can depend on both  $\epsilon$  and  $\zeta \sim 1 + r\sqrt{|K|}$

# Main results

$$\zeta \sim 1 + r\sqrt{|K|}$$

g-convex setting	Lower bound	Upper bound	Algorithm
(P1) Lipschitz, lo-dim		$O(\zeta d^2)$	Centerpoint method (Rusciano'19)
(P2) Lipschitz		$O\left(\frac{\zeta}{\epsilon^2}\right)$	Subgradient descent (Zhang & Sra'16)
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Curvature dependence in all upper bounds is unavoidable!



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Slight modification of hard function in C & Boumal'22

Volume of ball of radius  $r$ :  $\text{Vol} \sim e^{dr}$

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**Additive**

**Do not match!**


  
**Multiplicative**

Note:  $\kappa \geq \zeta$  and  $\frac{1}{\epsilon} \geq \zeta$

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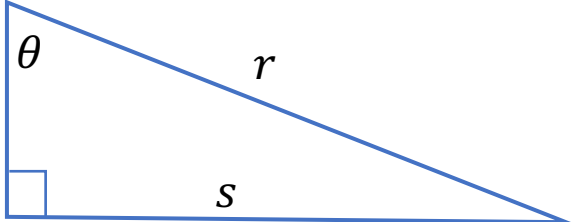
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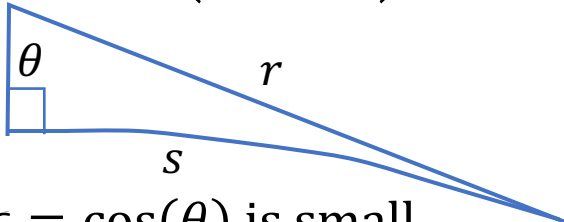
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(P3) Strongly g-convex			Centerpoint method (Rusciano'19) & Yang'22 / Luo et al.'22)
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Choosing halfspaces in right way yields lower bound  $\tilde{\Omega}\left(\frac{1}{\zeta^2 \epsilon^2}\right)$ .

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Smooth the Lipschitz lower bound with Moreau envelope (Guzman and Nemirovski'15):

$$f_\lambda(x) = \inf_{y \in \mathcal{M}} \left\{ f(y) + \frac{1}{2\lambda} \text{dist}^2(x, y) \right\}$$

# Smooth lower bound (P3)

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Surprising because no good notion of Fenchel dual on manifolds.

Yields lower bound  $\tilde{\Omega}\left(\frac{1}{\zeta\sqrt{\epsilon}}\right)$ .

# Tight lower bound for subgradient descent (P2)

New worst function in the world (*not* an extension of a Euclidean proof):

$$f(x) = \text{dist}(x, x^*) +$$



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$$f(x) = \text{dist}(x, x^*) + \max_i \left\{ \frac{1}{4\epsilon} \text{dist}(x, L_i) \right\}$$

$L_i$  are cleverly chosen hyperbolic halfspaces.

Yields  $\Omega\left(\frac{\zeta}{\epsilon^2}\right)$  lower bound for subgradient descent (with Polyak step size).

# Tight lower bound for subgradient descent (P2)

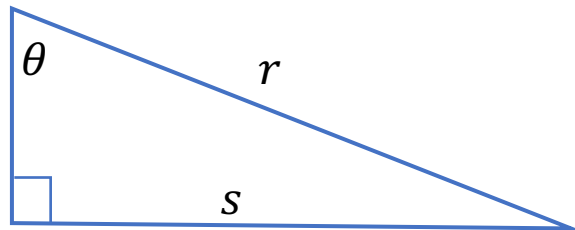
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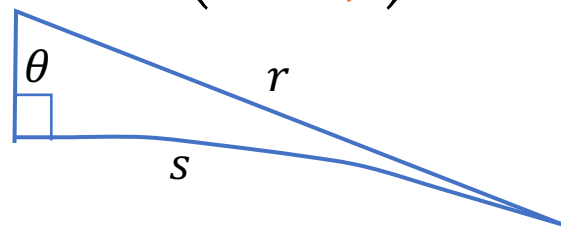
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$$s^2 = (1 - \epsilon^2)r^2$$



$$s^2 \approx \left(1 - \frac{\epsilon^2}{\zeta}\right)r^2$$



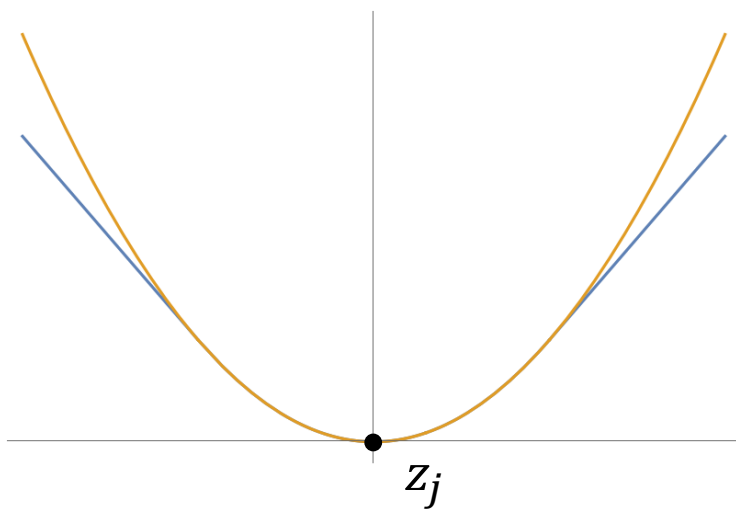
$\epsilon = \cos(\theta)$  is small

# Appendix

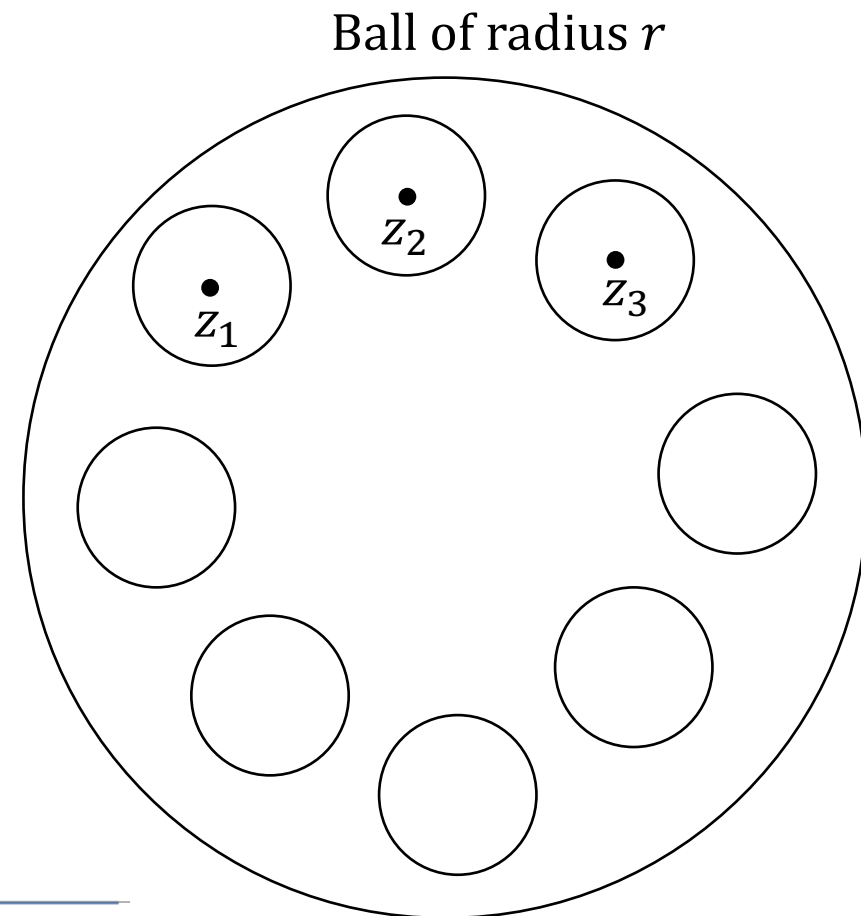
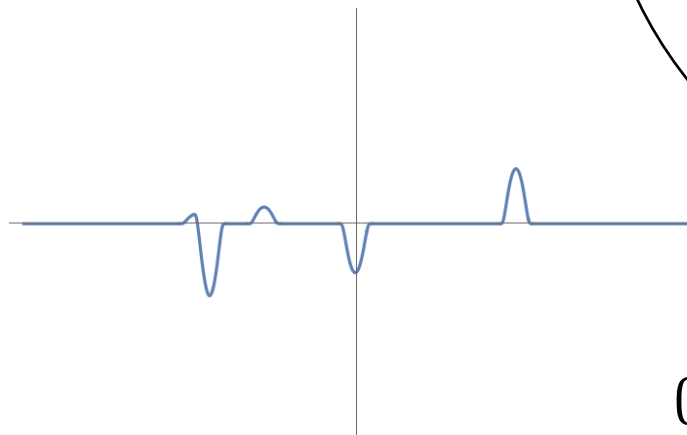
# Sec 3

Slight modification of hard function in C & Boumal'21

Volume of ball of radius  $r$ :  $\text{Vol} \sim e^{dr}$



+



Ball of radius  $\frac{r}{5}$   
(exponentially many of these)

# Lipschitz lower bound (P2) – Sec 4

In Euclidean space:  $f(x) = \max_{i=1,\dots,d} \{\langle s_i e_i, x \rangle\}$ ,  $s_i \in \{-1, +1\}$

Fundamental difficulty: No good notion of linear functions on manifolds.

- Every convex function equals a **sup of affine functions**, which are **themselves convex**.
- Every g-convex function equals a **sup of  $x \mapsto F_i + \langle g_i, \log_{y_i}(x) \rangle$** , but these are **not g-convex**.

Totally geodesic submanifolds:  $S_i = \{x \in \mathcal{M} = \mathbb{H}^d : \langle g_i, \log_{y_i}(x) \rangle = 0\}$

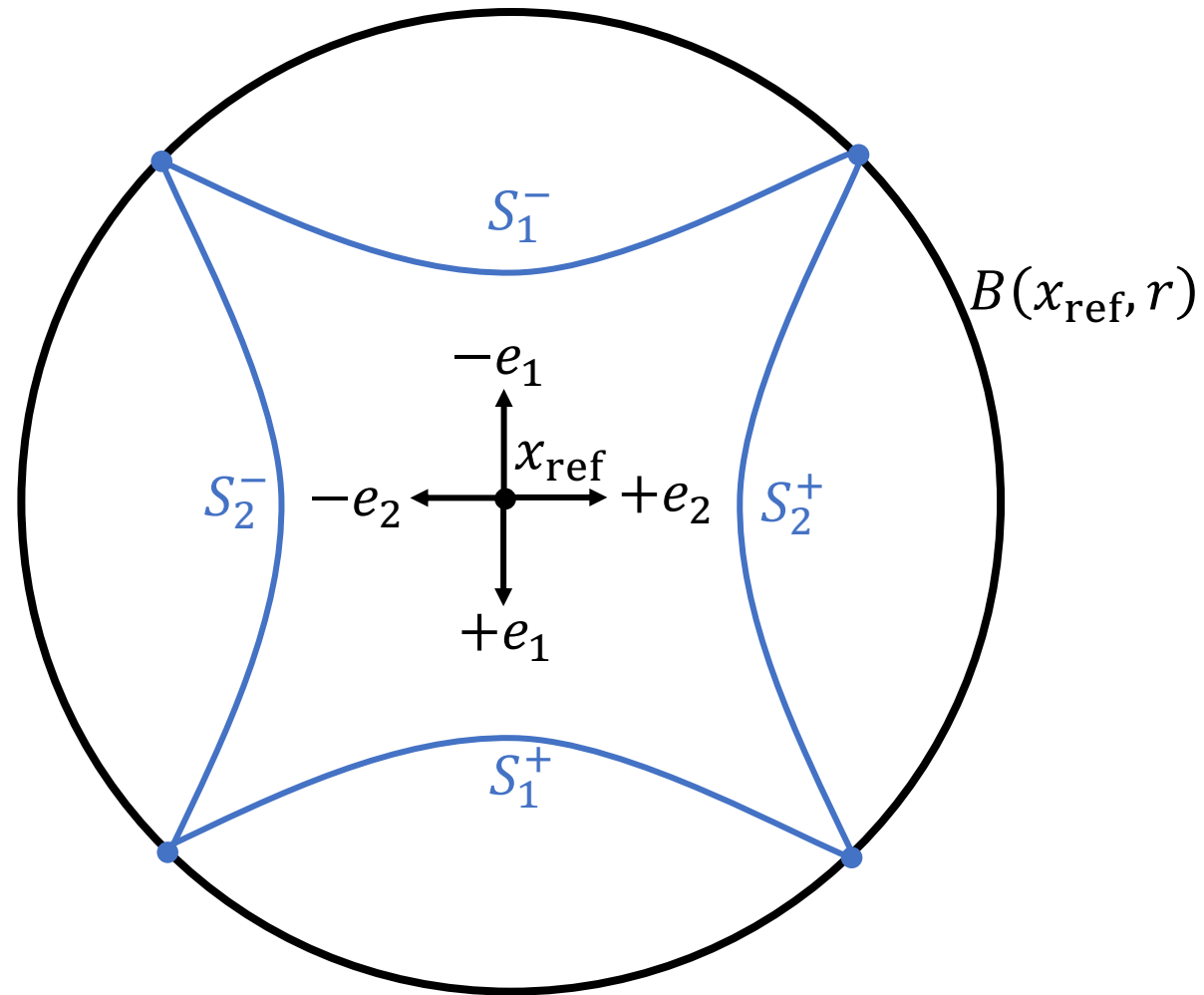
- Generalization of affine subspace

Idea: use **max of distance to totally geodesic submanifolds**:

$$f(x) = \max_{i=1,\dots,d} \{\text{dist}(x, S_i^{s_i})\}, s_i \in \{-1, +1\}$$

Choosing halfspaces in right way yields lower bound  $\tilde{\Omega}\left(\frac{1}{\zeta^2 \epsilon^2}\right)$ .

# Lipschitz lower bound (P2) – Sec 4



# Smooth lower bound (P3) – Sec 5

~~Worst function in world:  $f(x) = x_{(1)}^2 - 2x_{(1)} + \sum_i (x_{(i)} - x_{(i+1)})^2 + x_{(n)}^2$~~

Smooth Lipschitz lower bound with Moreau envelope (Guzman and Nemirovski'15):

$$f_\lambda(x) = \inf_{y \in \mathcal{M}} \left\{ f(y) + \frac{1}{2\lambda} \text{dist}^2(x, y) \right\}$$

Surprising because no good notion of Fenchel dual on manifolds.

Yields lower bound  $\tilde{\Omega}\left(\frac{1}{\zeta\sqrt{\epsilon}}\right)$ .

# Tight lower bound for subgradient descent (P2)

## – Sec 6

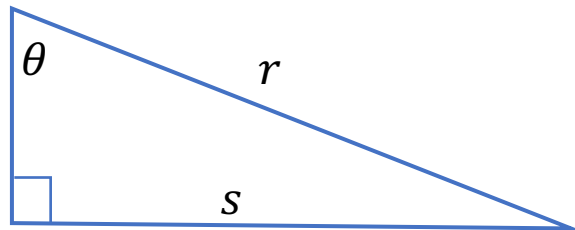
New worst function in the world (not an extension of a Euclidean proof):

$$f(x) = \text{dist}(x, x^*) + \max_{i=0, \dots, d-2} \left\{ \frac{1}{4\epsilon} \text{dist}(x, L_i) \right\}$$

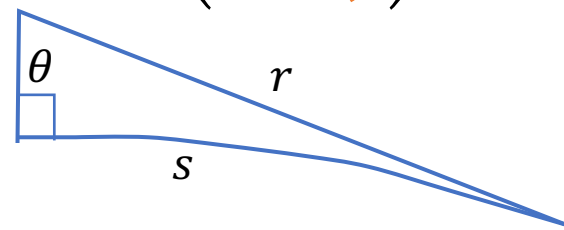
$L_i$  are cleverly chosen hyperbolic halfspaces.

Yields  $\Omega\left(\frac{\zeta}{\epsilon^2}\right)$  lower bound for subgradient descent (with Polyak step size).

$$s^2 = (1 - \epsilon^2)r^2$$



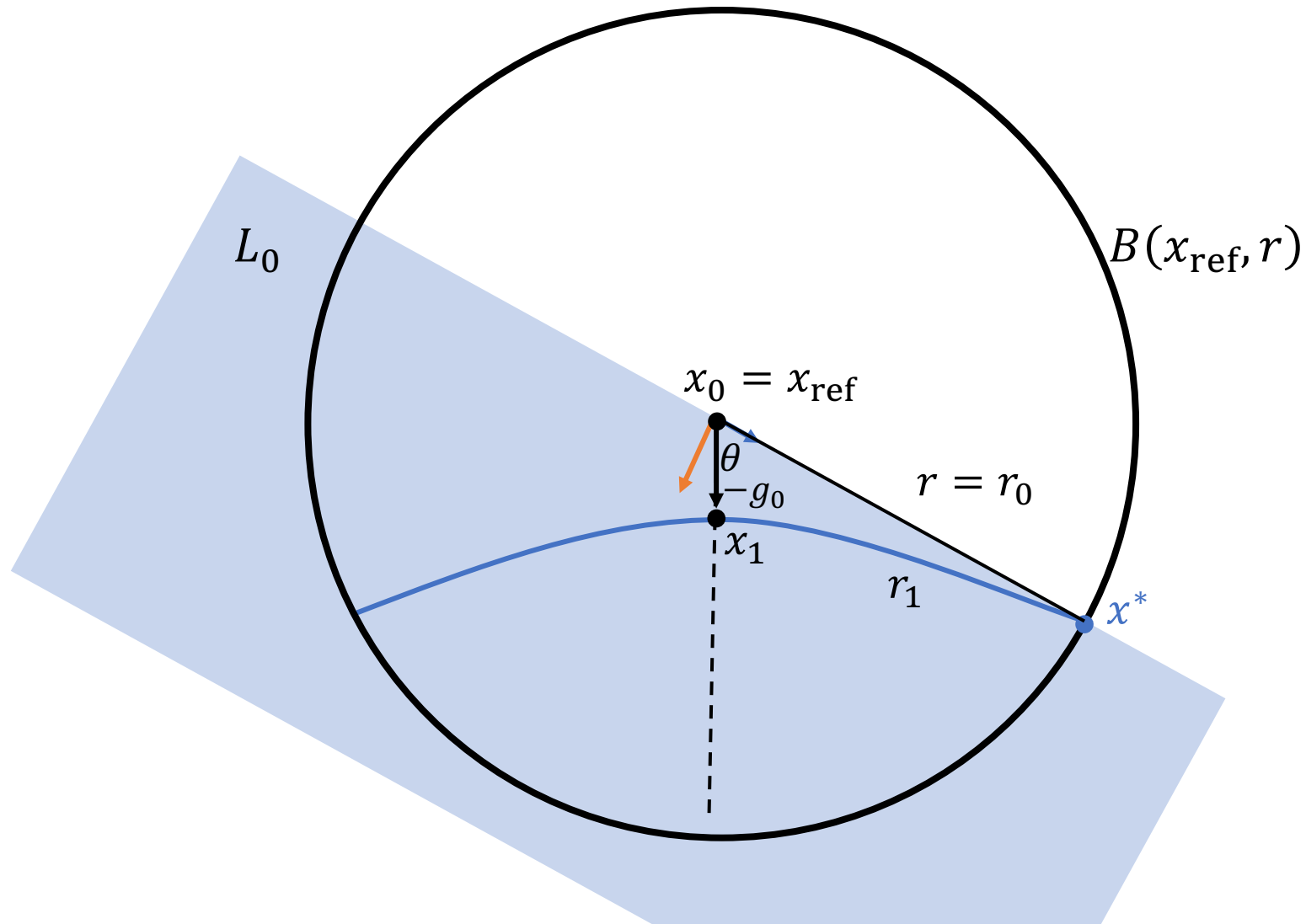
$$s^2 \approx \left(1 - \frac{\epsilon^2}{\zeta}\right)r^2$$



$\epsilon = \cos(\theta)$  is small



# Tight lower bound for subgradient descent (P2) – Sec 6



# Cutting-planes game (proxy for (P1)) – Sec 7

Cutting planes game:

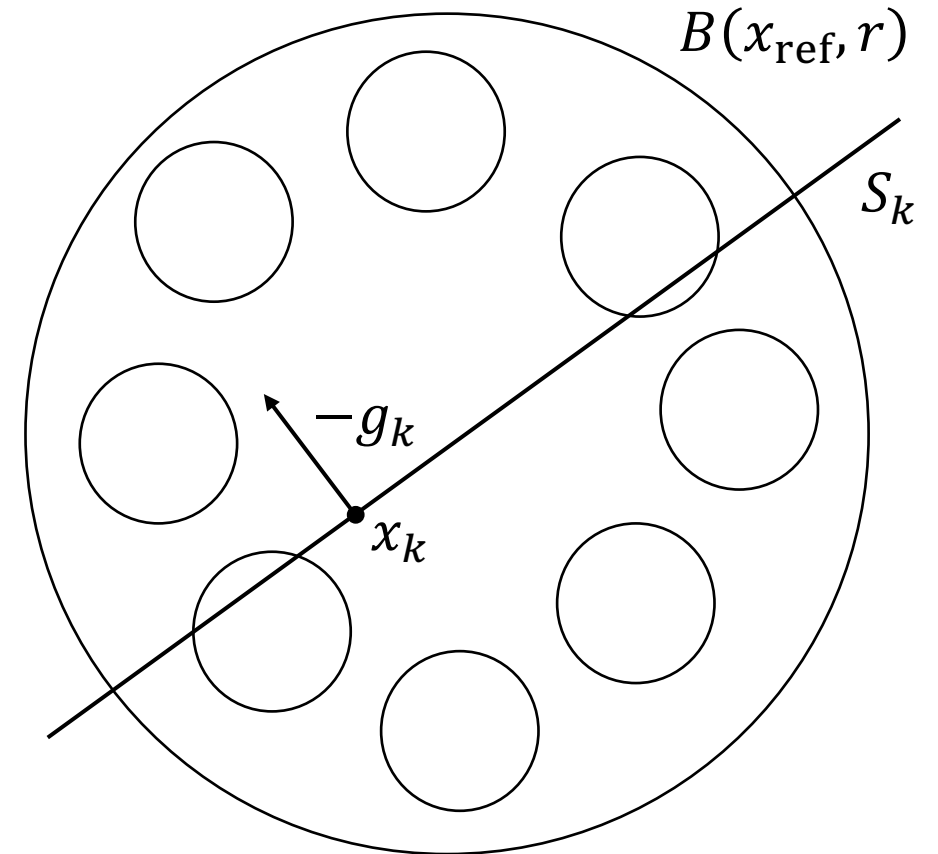
- Initially given  $B(x_{\text{ref}}, r)$  containing  $x^*$
- Upon query  $x_k$ , oracle returns tangent vector  $g_k$  such that  $\langle g_k, \log_{x_k}(x^*) \rangle \leq 0$
- Task: Find  $x_k$  so that  $B(x^*, \epsilon r)$  intersects  $S_k = \{x \in \mathcal{M} : \langle g_k, \log_{x_k}(x) \rangle = 0\}$

Usual lower bounds proof (tessellation of cube) does not generalize (Nesterov'04).

**Width-bounded separators** (Kisfaludi-Back'20, Fu'11):

- Given a collect of balls and a point  $x_k$ , there exists a totally geodesic submanifold  $S_k$  through  $x_k$  which intersects a small fraction of these balls

Yields  $\tilde{\Omega}(\zeta d)$  lower bound for cutting-plane schemes, e.g., center-of-gravity method.



# G-convex interpolation – Sec 8

Interpolation is crucial for building lower bounds (Taylor et al.'16).

A collection of function values and tangent vectors  $(F_i, x_i, g_i)_{i=1}^N$  is interpolated by a g-convex function  $f$  if  $f(x_i) = F_i$  and  $g_i \in \partial f(x_i)$  for all  $i$ .

A **convex** function interpolates  $(F_i, x_i, g_i)_{i=1}^N$  **if and only if**

$$F_j \geq F_i + \langle g_i, x_j - x_i \rangle \text{ for all } i, j$$

For **g-convex** functions the analogous **naïve necessary conditions are *not* sufficient for interpolation even** for just 3 points:

- There exists  $(F_i, x_i, g_i)_{i=1}^3$  such that  $F_j \geq F_i + \langle g_i, \log_{x_j}(x_j) \rangle$  for all  $i, j$ , yet this data cannot be interpolated by a g-convex function.